Searching for Optimal Strategies in Proof-of-Stake Mining Games with Access to External Randomness

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Overview

- 1. Motivation
- 2. Game
- 3. Prior Work
- 4. Structured Strategies
- 5. Symmetrical States
- 6. Non-Checkpoint Finality
- 7. n-Deficit Tolerance Family of Strategies
- 8. Automating this Search
- 9. Conclusion

Proof-of-Work (PoW) Mining Protocol

Proof-of-Work

- 1. Present a string s for which H(s) < target.
- 2. Mine a block of coins.



H is non-invertible, so *s* is a **needle in a haystack**.

Computing hashes over random strings consumes electricity.

Proof-of-Work (PoW) Mining Protocol

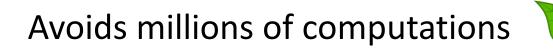
How much electricity does proof-of-work consume?



Proof-of-Stake (PoS) Mining Protocol

Proof-of-Stake (w / External Randomness)

- 1. External source chooses $x \in \{1, \text{total number of coins}\}$
- 2. If you own coin *x*, mine a block of coins.





Project Goal

Investigate whether the PoS mining protocol is a viable alternative to the PoW mining protocol

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Cryptocurrency Mining Game

Use a **2-player** game to model mining cryptocurrency under PoS.

Attacker has strength (probability of mining a block) α . Defender has strength (probability of mining a block) $1 - \alpha$.



Attacker hidden blocks:



Current timestep is bolded.

Cryptocurrency Mining Game: Demo

Attacker hidden blocks:

-	
Timestep	Miner
0	N/A

Current timestep is bolded.

Cryptocurrency Mining Game: Demo

Timestep Miner 0 N/A

Blockchain (Longest path marked with heavy arrows.)

Attacker hidden blocks:

Blockchain

0

Current timestep is bolded.

Attacker hidden blocks:	Timestep	Miner
	0	N/A
Defender hidden blocks:		
(Longest path marked with heavy arrows.)	<u>Game Transcrip</u>	<u>t:</u>
	Game setup.	

Current timestep is bolded.

X	
Timestep	Miner
0	N/A
1	D

Blockchain (Longest path marked with heavy arrows.)

Attacker hidden blocks:

Defender hidden blocks:

Game Transcript:

Defender mines block 1.



Current timestep is bolded.

Cryptocurrency Mining Game: Demo

TimestepMiner0N/A1D

Blockchain (Longest path marked with heavy arrows.)

0

Attacker hidden blocks:

Defender hidden blocks:

Game Transcript:

Defender adds block 1 to the blockchain, pointing to 0.

15

Current timestep is bolded.

Cryptocurrency Mining Game: Demo

2

TimestepMiner0N/A1D2A

Blockchain (Longest path marked with heavy arrows.)

Attacker hidden blocks:

Defender hidden blocks:

Game Transcript:

Attacker mines block 2.



16

Current timestep is bolded.

Cryptocurrency Mining Game: Demo

2

3

TimestepMiner0N/A1D2A3A

Blockchain (Longest path marked with heavy arrows.)

Attacker hidden blocks:

Defender hidden blocks:

Game Transcript:

Attacker mines block 3.



Current timestep is bolded.

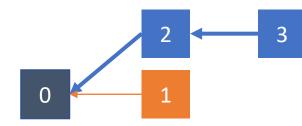
Cryptocurrency Mining Game: Demo

TimestepMiner0N/A1D2A3A

Blockchain (Longest path marked with heavy arrows.)

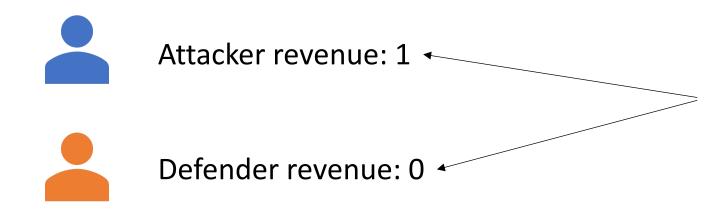
Attacker hidden blocks:

Defender hidden blocks:

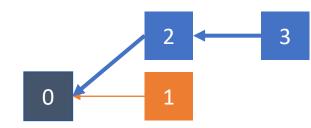


Game Transcript:

Attacker adds blocks 2 and 3 to the blockchain, pointing 3 to 2 and 2 to 0.



Proportion of blocks they own in the **longest path** in the blockchain.



Honest Mining

Players are *supposed* to use the **honest mining strategy** (HONEST):

If you mine a block, publish it on the longest chain.

Selfish Mining

Players can cheat and do better than HONEST.

Such **selfish mining** threatens the adoption of PoS.



Robustness of PoS to Attack

Let **robustness** be the minimum strength α necessary to conduct an attack.

Denote this as α^{PoS} .

Project Goal

Investigate whether the PoS mining protocol is a viable alternative to the PoW mining protocol ...

Project Goal

Investigate whether the PoS mining protocol is a viable alternative to the PoW mining protocol ...

by bounding α^{PoS} , the robustness of the PoS mining protocol to an attack.



Attacker hidden blocks:



Defender hidden blocks:





Attacker hidden blocks:





Defender hidden blocks:





Attacker hidden blocks:





Defender hidden blocks:











Defender hidden blocks:





Attacker hidden blocks:





Defender hidden blocks:











Defender hidden blocks:

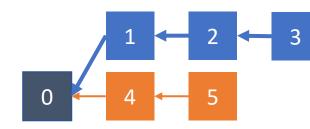




Attacker hidden blocks:



Defender hidden blocks:

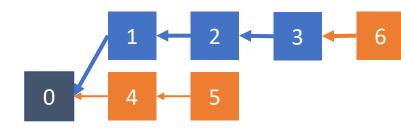




Attacker hidden blocks:



Defender hidden blocks:



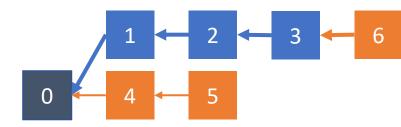


Attacker hidden blocks:





Defender hidden blocks:

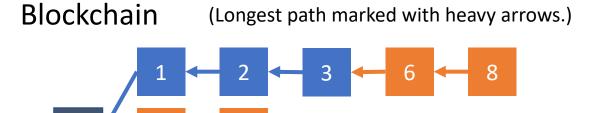




Attacker hidden blocks:





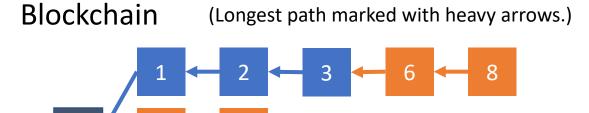




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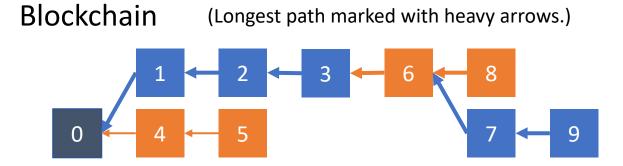






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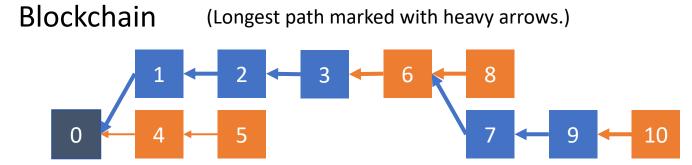






Attacker hidden blocks:





Overview

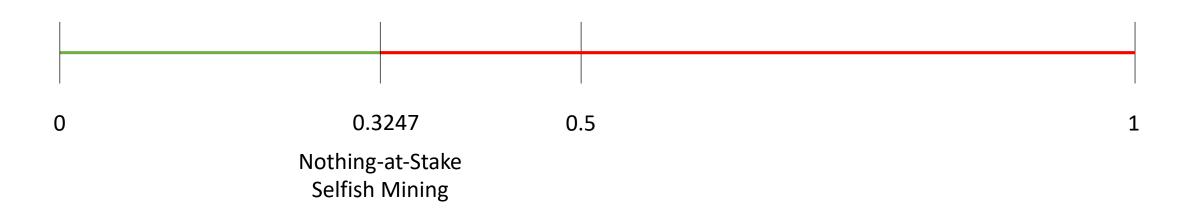
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 $0 < \alpha^{PoS} < 1$

0

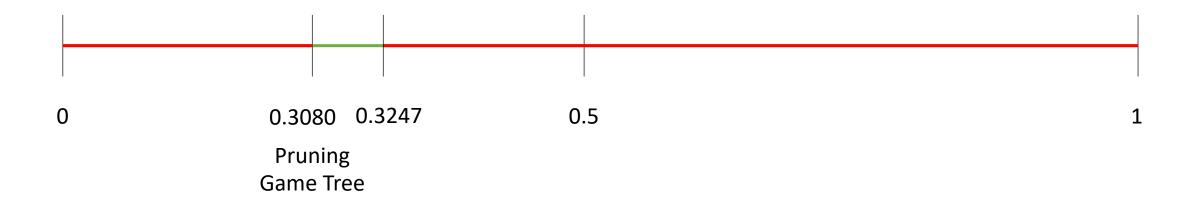
1

$0 \le \alpha^{PoS} \le 0.3247$

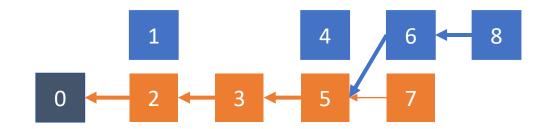


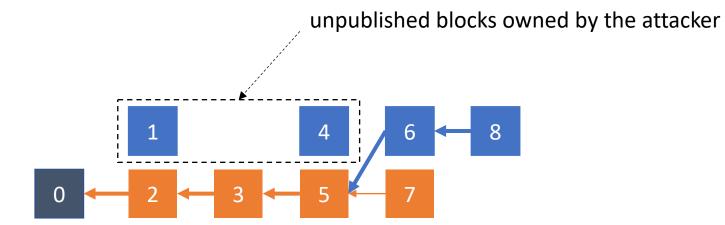
* Not drawn to scale.

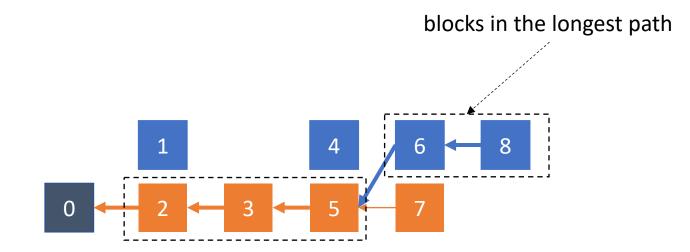
$0.3080 \le \alpha^{PoS} \le 0.3247$

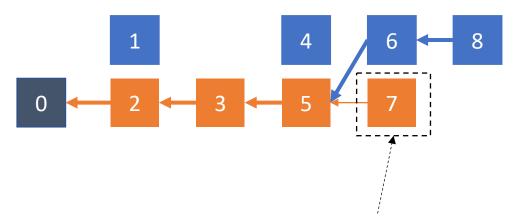


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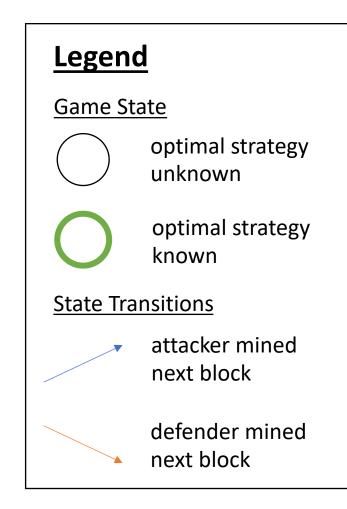




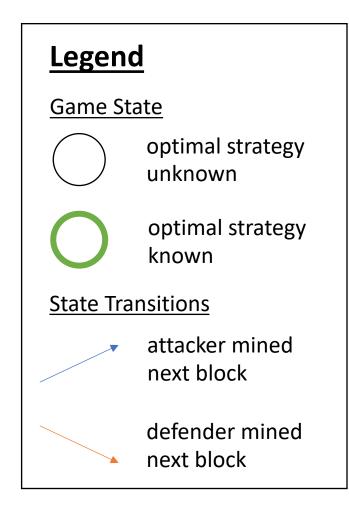


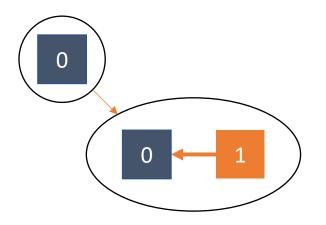


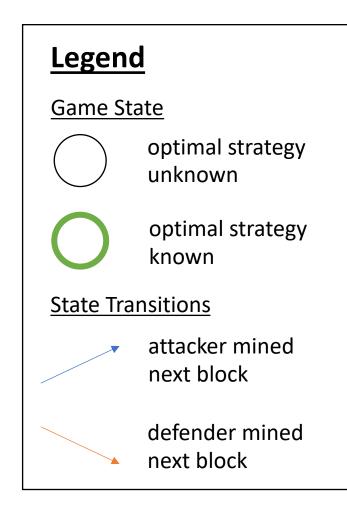
orphaned block; block forked from the longest path

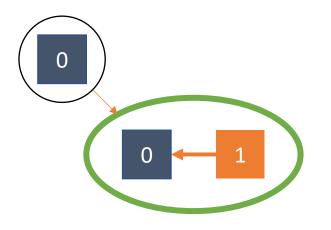


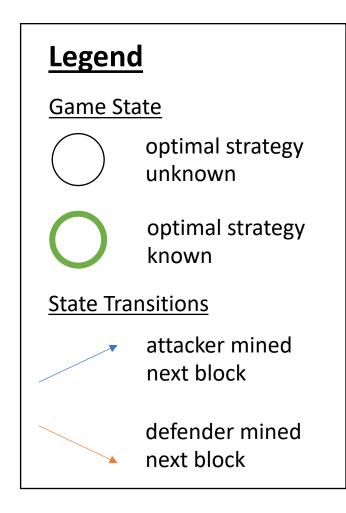


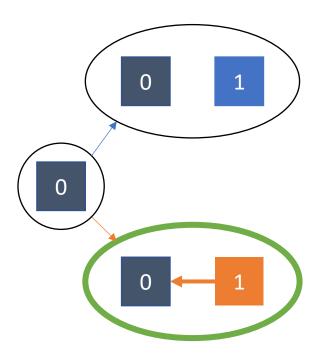


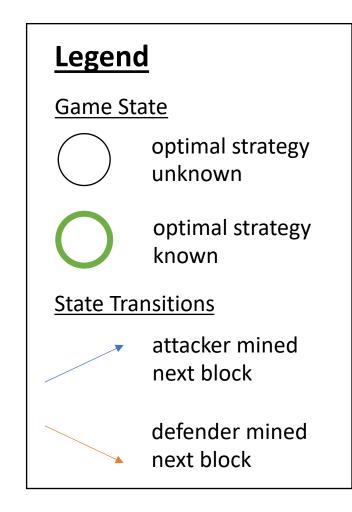


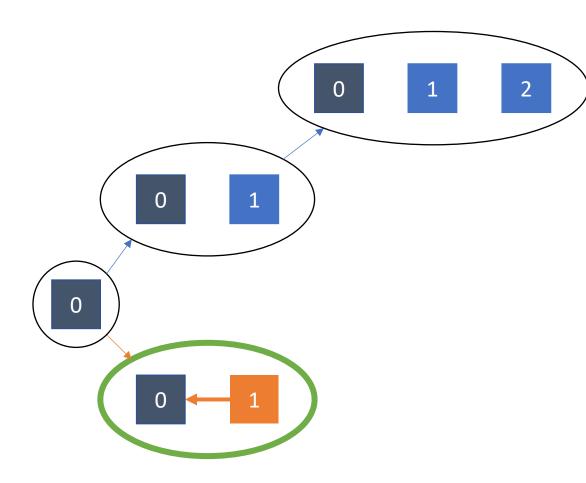


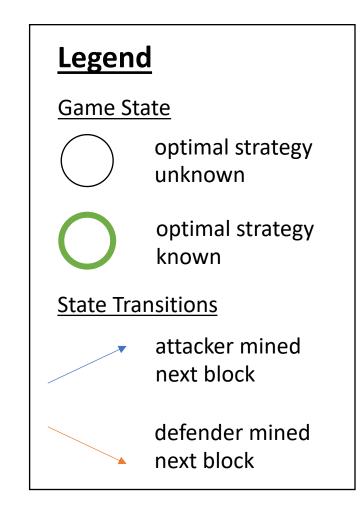


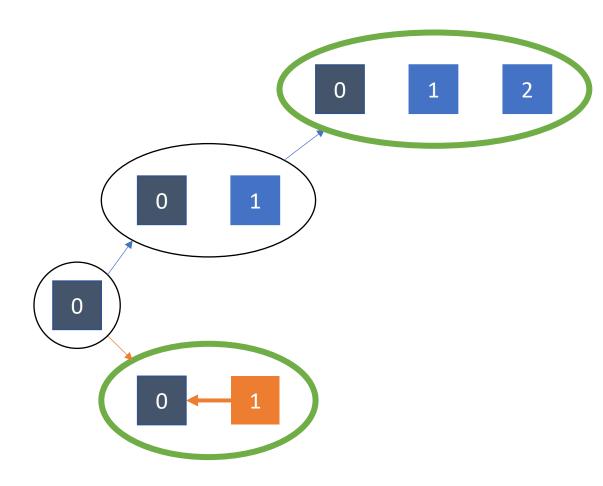


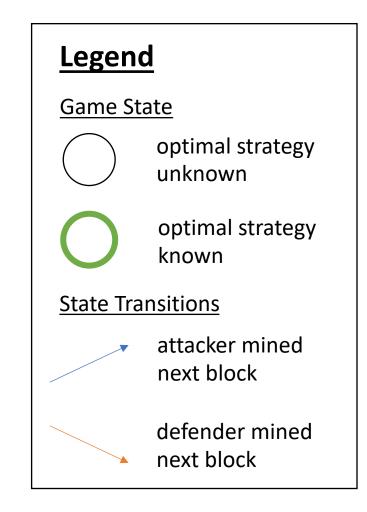


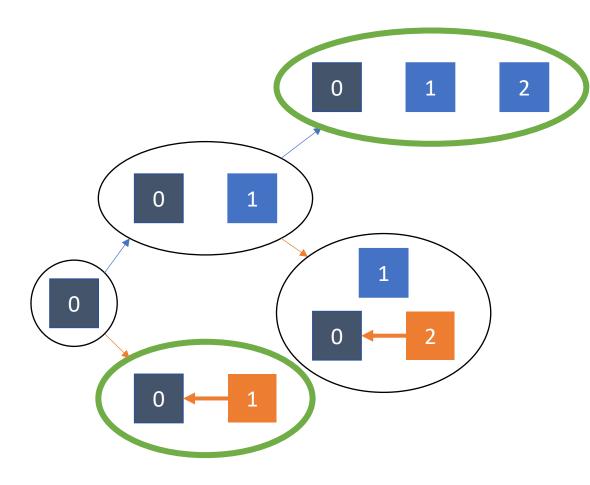


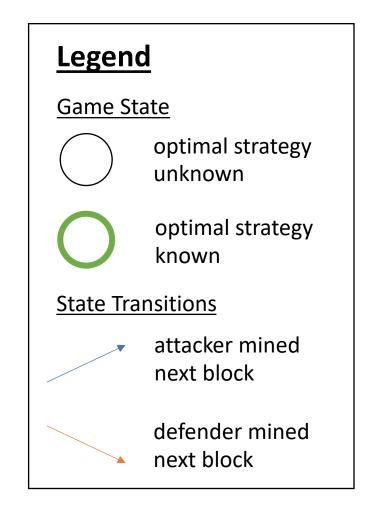


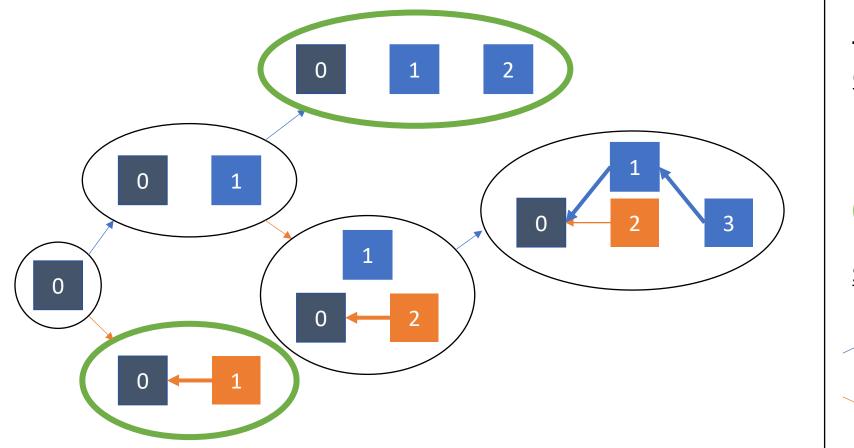


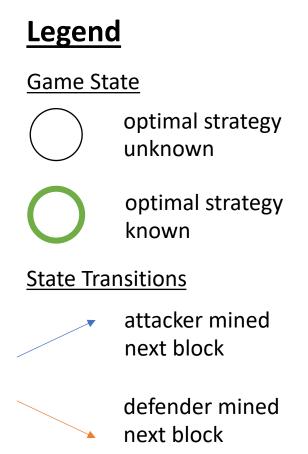


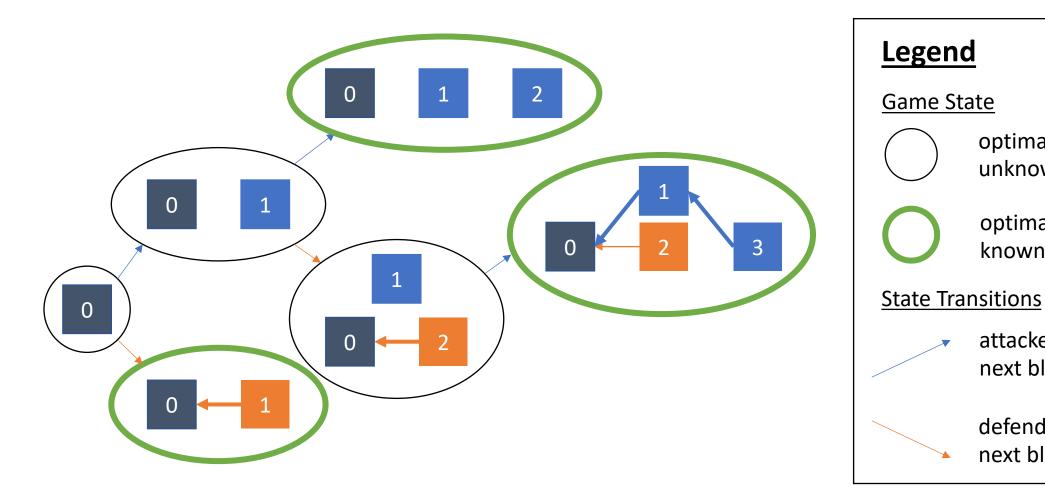












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optimal strategy

optimal strategy

attacker mined

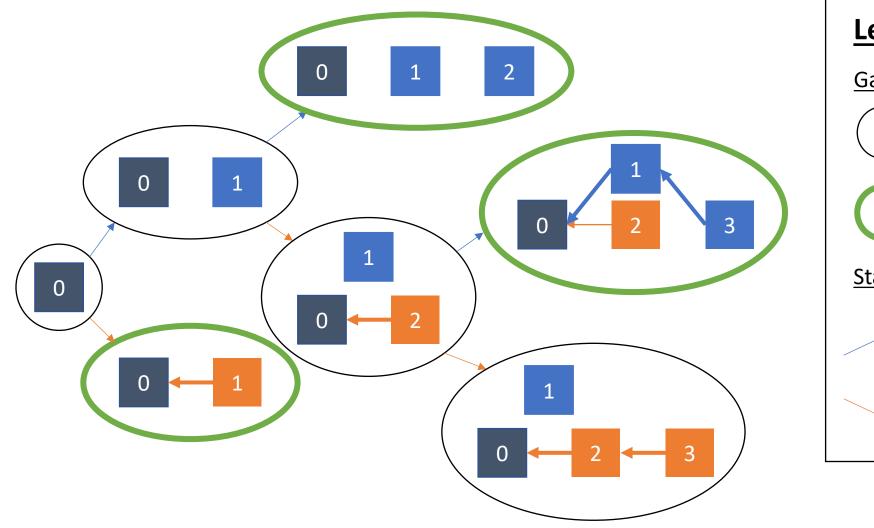
defender mined

next block

next block

unknown

known



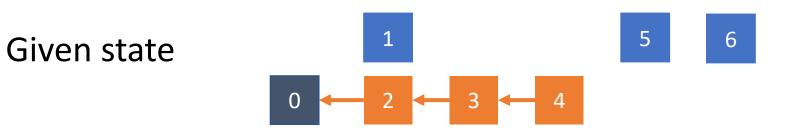
Legend Game State optimal strategy unknown optimal strategy known **State Transitions** attacker mined next block defender mined next block

Overview

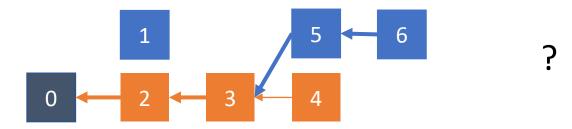
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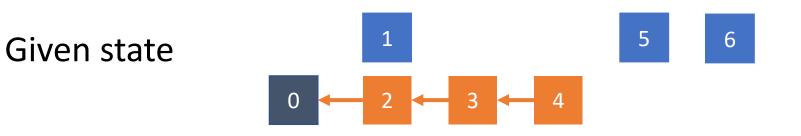
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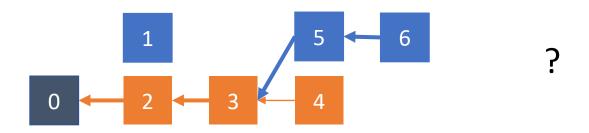


should a strategy take action

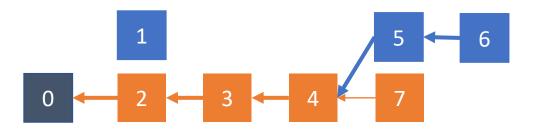


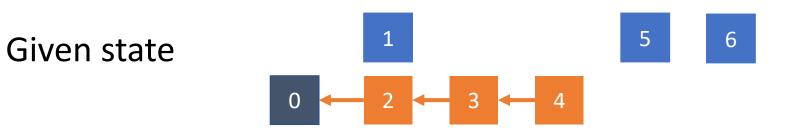


should a strategy take action

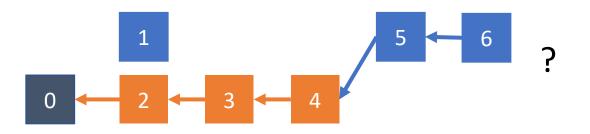


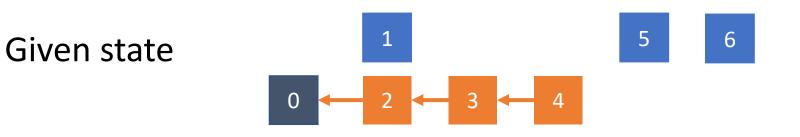
NO! Just wait... can still publish even in the worstcase scenario that defender mines the next block.



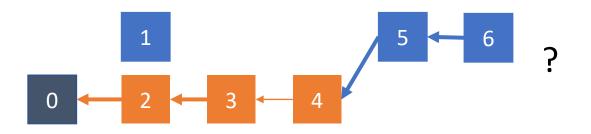


should a strategy take action

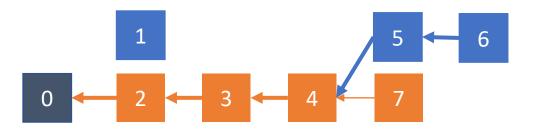


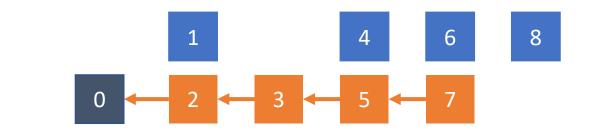


should a strategy take action



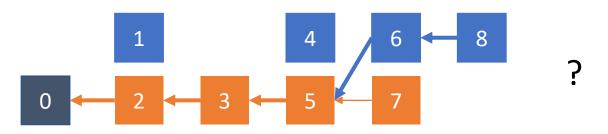
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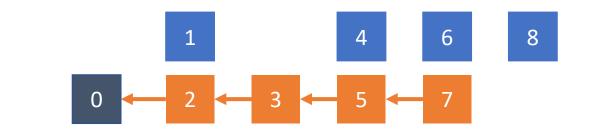




should a strategy take action

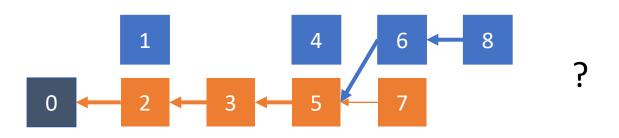
Given state



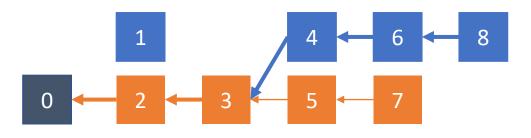


should a strategy take action

Given state



NO! Why wouldn't the strategy also publish block 4?



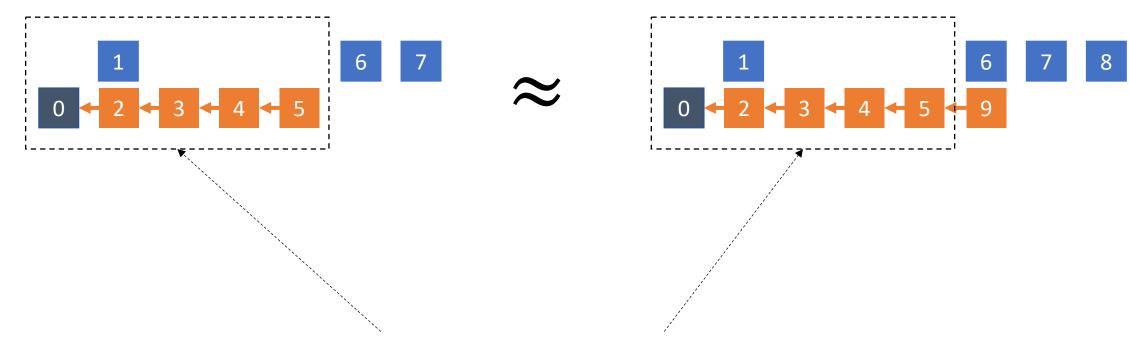
A *structured strategy* obeys these and several other "intuitive" properties so there are fewer actions to compare at any state.

Without loss of generality, an optimal strategy is *structured*.

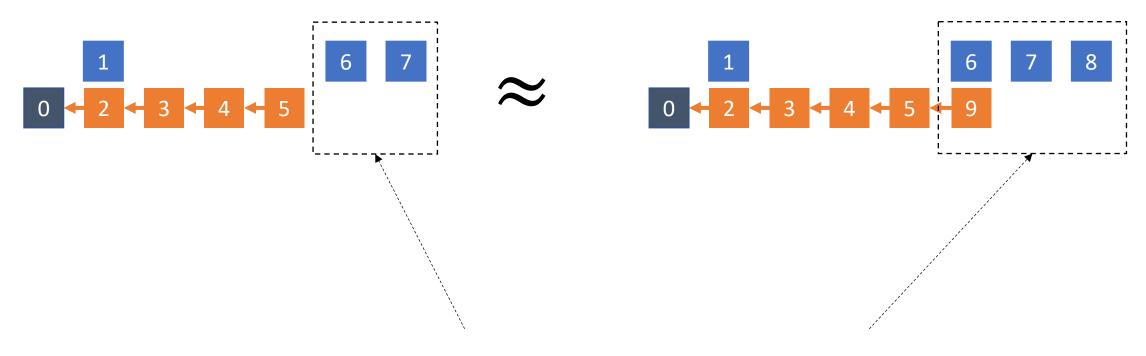
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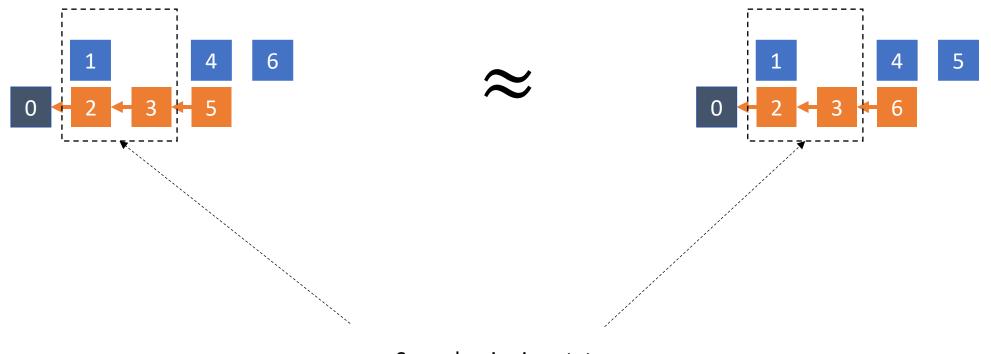
Same beginning state.



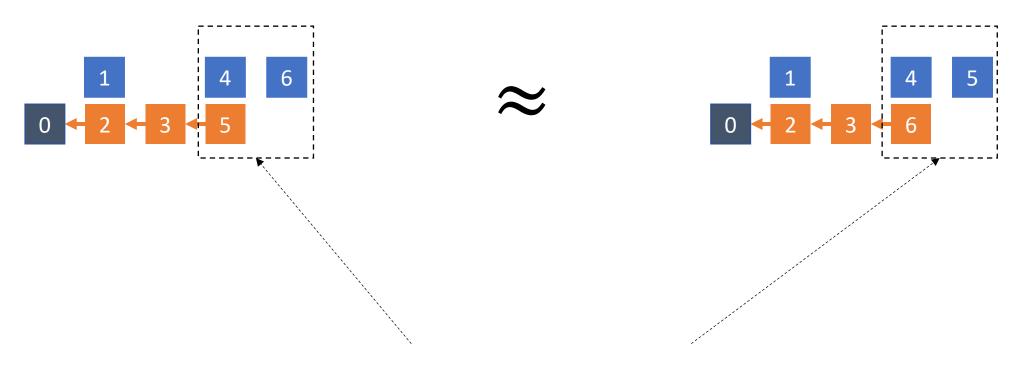
Same lead over the subsequent block.



The values of these states are related by *exactly* the difference in the number of attacker blocks!



Same beginning state.



Only difference is 6 can never be published on 5, but why would they want to do that anyways?



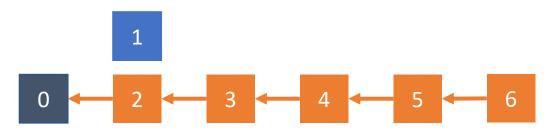
The values of these states are *exactly* equal!

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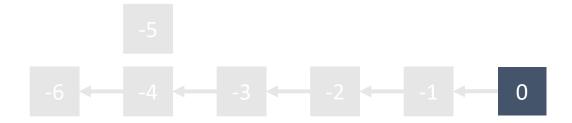
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Non-Checkpoint Finality

If the game reaches

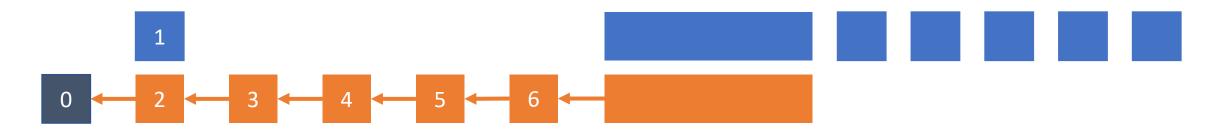


then the optimal strategy is to just "reset" the game, or



Non-Checkpoint Finality

Proof Sketch: Suppose that the attacker ever has the chance to publish block 1 such that it enters the longest path. Then the game state looks something like



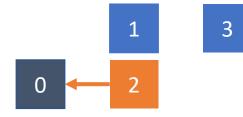
That is, the attacker has a lead of five blocks.

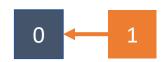
But, a lead of five blocks let's you do things which are *strictly better* than publishing block 1.

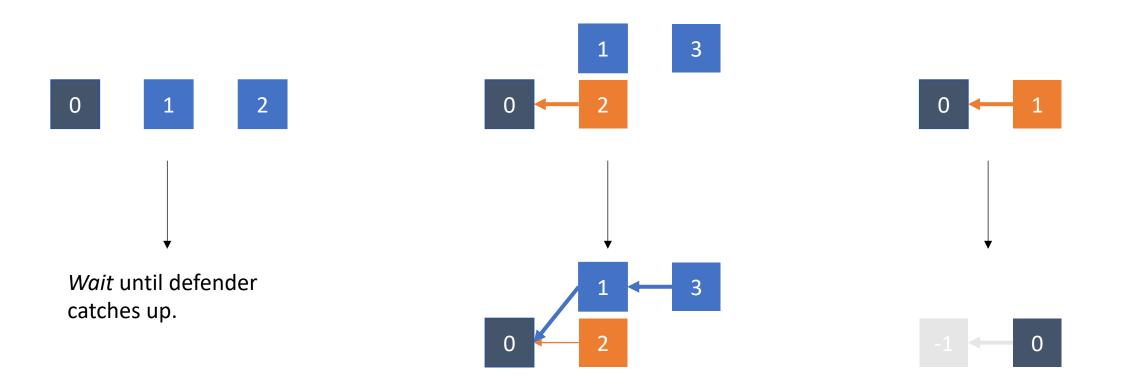
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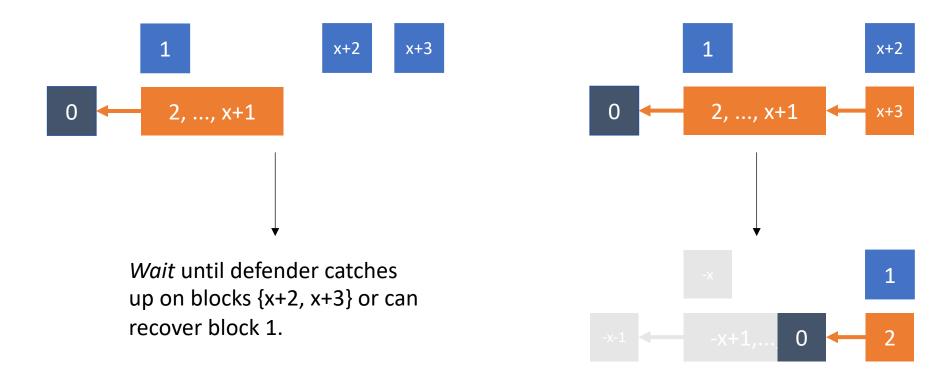


- Take optimal action at states where this is known.
- Take *reasonable* actions at other states that are interesting.

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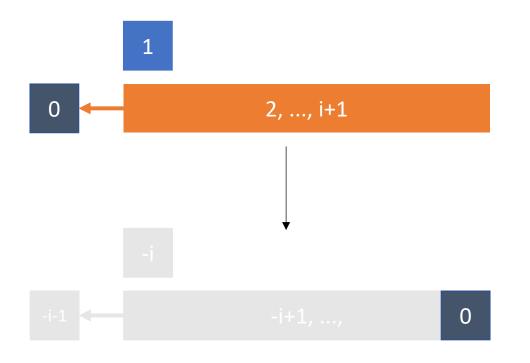
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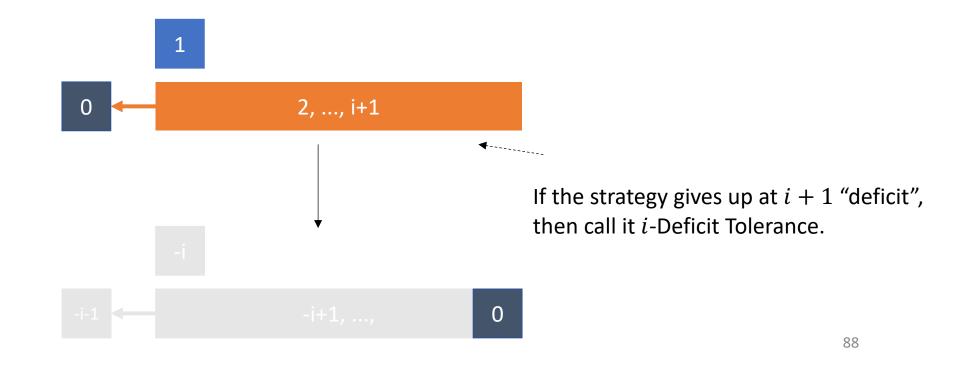
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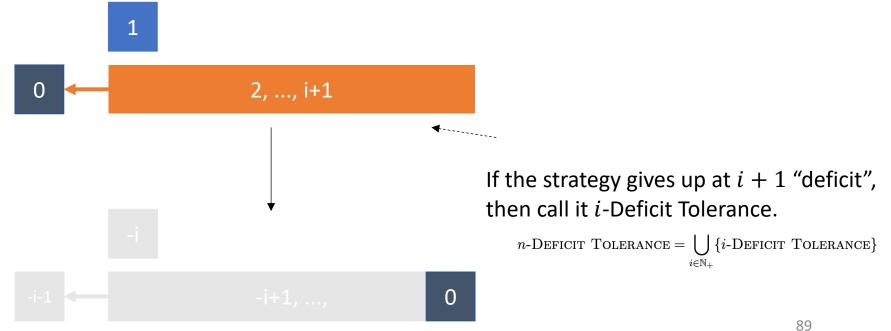
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- Take optimal action at states where this is known.
- Take *reasonable* actions at other states that are interesting.
- Wait elsewhere.

Strategy π	$\min\{\alpha \in [0,1] \mid \operatorname{Rev}(\pi,\alpha) > \operatorname{Rev}(\operatorname{Honest},\alpha)\}$
1-Deficit Tolerance (SM)	0.333333
2-Deficit Tolerance (NSM)	0.324718
3-Deficit Tolerance	0.323577
4-Deficit Tolerance	0.323489
5-Deficit Tolerance	0.323534
6-Deficit Tolerance	0.323572

Smallest mining strength where it outperforms the honest strategy.

Strategy π	$\min\{\alpha \in [0,1] \mid \operatorname{Rev}(\pi,\alpha) > \operatorname{Rev}(\operatorname{HONEST},\alpha)\}$
1-Deficit Tolerance (SM)	0.333333
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Prior

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	3-Deficit Tolerance	0.323577
· 🤞 🛛	4-Deficit Tolerance	0.323489
1	5-Deficit Tolerance	0.323534
	6-Deficit Tolerance	0.323572

 $0.3080 \leq \alpha^{PoS} \leq 0.3247$

0.3080 Pruning Game Tree 0.3247

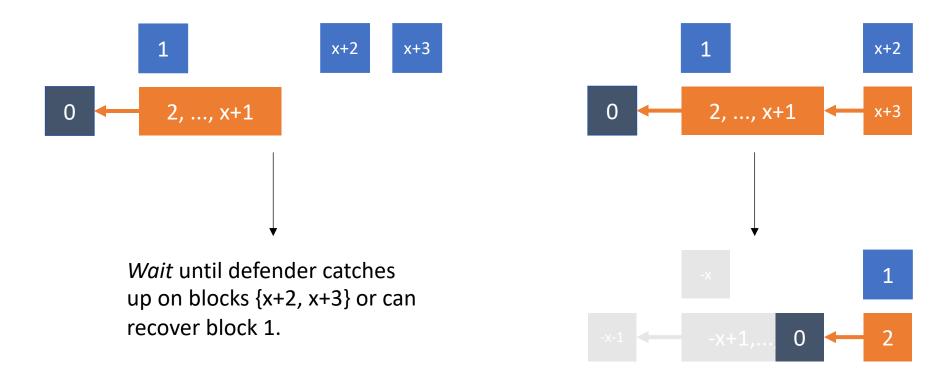
Nothing-at-Stake Selfish Mining

$0.3080 \leq \alpha^{PoS} \leq 0.3235$

0.3080	0.3235	0.3247
Pruning Game Tree	4-Deficit Tolerance	Nothing-at-Stake Selfish Mining

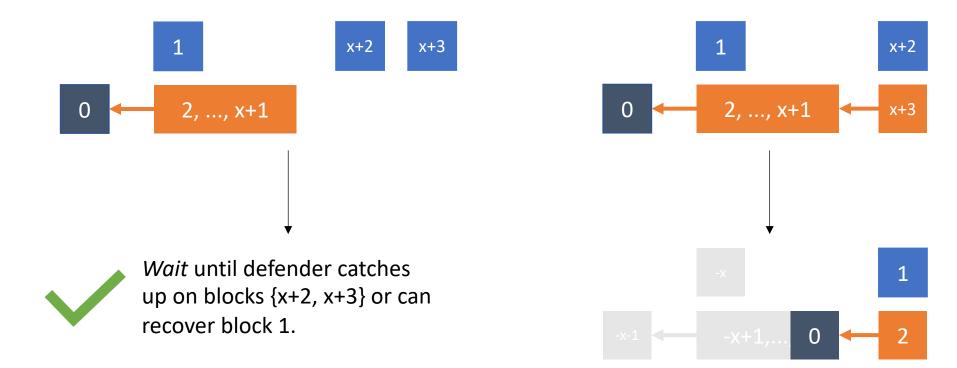
• Take optimal action at states where this is known.

evisited



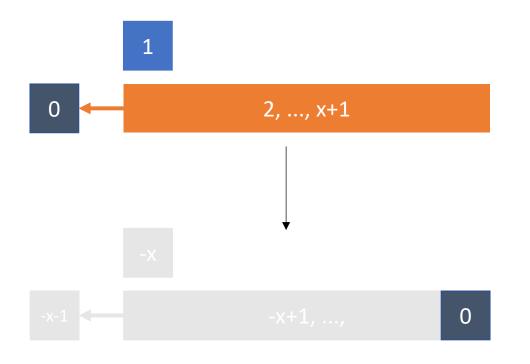
• Take optimal action at states where this is known.

evisited



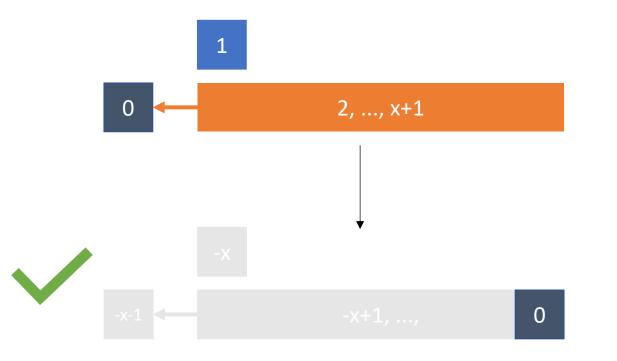
• Take optimal action at states where this is known.

revisited



• Take optimal action at states where this is known.

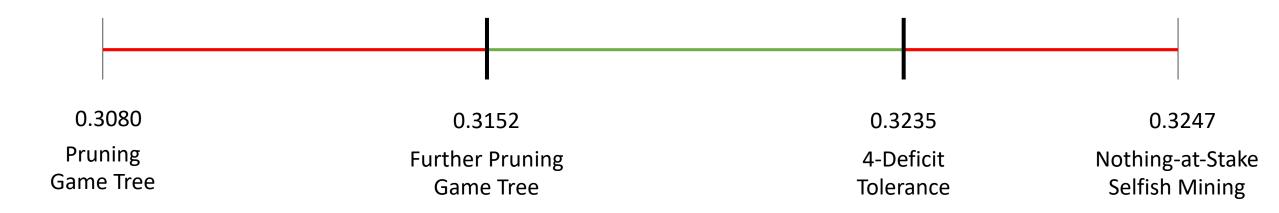
revisited



$0.3080 \leq \alpha^{PoS} \leq 0.3235$

0.3080	0.3235	0.3247
Pruning Game Tree	4-Deficit Tolerance	Nothing-at-Stake Selfish Mining

$0.3152 \leq \alpha^{PoS} \leq 0.3235$



Overview

- 1. Motivation
- 2. Game
- 3. Prior Work
- 4. Structured Strategies
- 5. Symmetrical States
- 6. Non-Checkpoint Finality
- 7. n-Deficit Tolerance Family of Strategies
- 8. Automating this Search
- 9. Conclusion

Automating this Search

<u>Algorithm</u>:

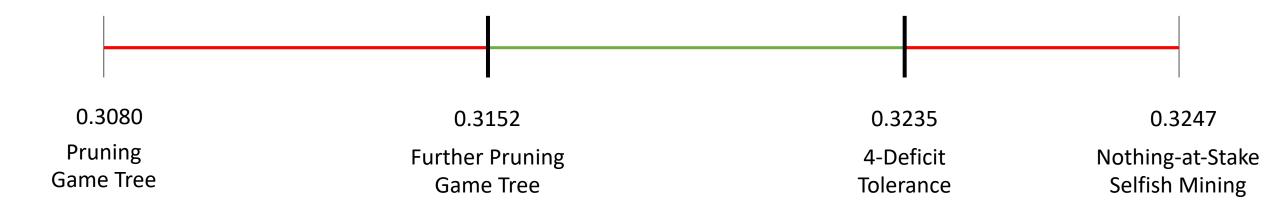
- 1. Simulate all reachable states with at most *n* blocks.
- 2. Lower and upper bound the value of all states with *n* blocks.
- 3. For i = n 1, ..., 0:
 - a. Lower and upper bound the value of all states with *i* blocks.

The above results on structured strategies, symmetrical states, and non-checkpoint finality, make this *computationally feasible*.

https://thesis.anthonyhein.com/materials/code-results/index.html

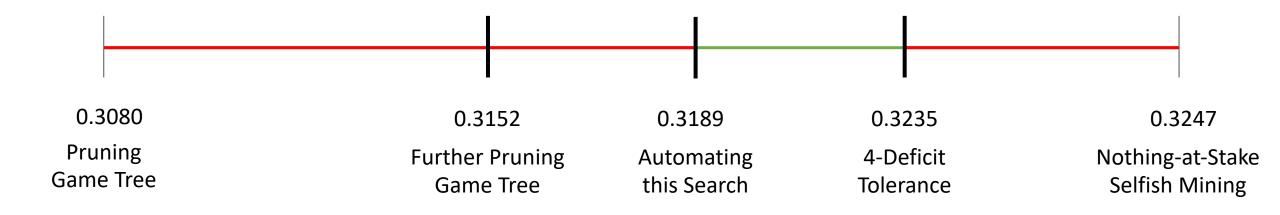
Automating this Search

$0.3152 \leq \alpha^{PoS} \leq 0.3235$



Automating this Search

$0.3189 \leq \alpha^{PoS} \leq 0.3235$



Overview

- 1. Motivation
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- 6. Non-Checkpoint Finality
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- 8. Automating this Search
- 9. Conclusion

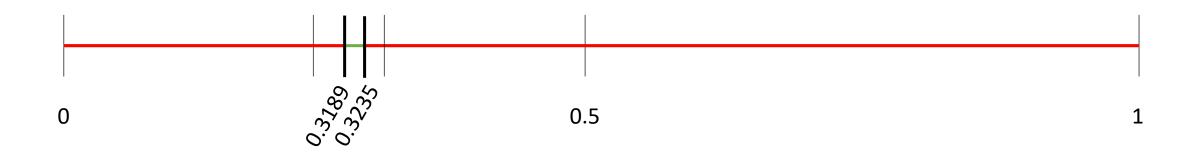
Conclusion

$0.3080 \leq \alpha^{PoS} \leq 0.3247$



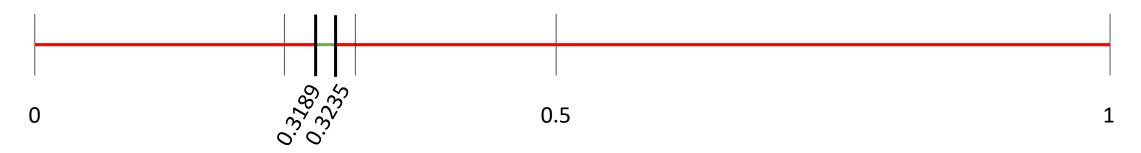
Conclusion

$0.3189 \leq \alpha^{PoS} \leq 0.3235$



Conclusion

Assuming a Bitcoin-like cryptocurrency, increasing mining strength by 0.001 costs \$900,000,000.



Acknowledgements

Professor Matt Weinberg Doctor Matheus V. X. Ferreira Professor Mark Braverman Briana Macedo



All materials can be found at https://thesis.anthonyhein.com

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<u>Password</u>: pledge-my-honor

Questions?

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